

Exploring Efficient Microservice Level Parallelism

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Shanghai Jiao Tong University

- **1. Introduction: Monolithic to Microservice**
- 2. Motivation: Microservice Characterization
- 3. Microservice Level Parallelism
- 4. Evaluation: Effectiveness, Efficiency, Performance
- 5. Conclusions and Future Work



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The Emerging Microservice Architecture

Microservice disaggregates a monolithic application into many tiny services, each of whom can be managed and tested independently.



Monolithic Architecture

Microservice-based Architecture

The Revolution of Cloud Services

Microservices empower organizations to build and run scalable, agile, and resilient applications in dynamic environments.



Industrial Applications of Microservices

Many IT companies such as Alibaba and Amazon are actively embracing this

new software development paradigm.



Prior Work on Microservices

Prior work: microservice design and optimization.



We enhance the QoS-aware performance optimization for microservices.

> Fully exploiting the unique characteristics of microservices.

Different Level of Parallelisms

Parallelism has been exploited at various levels of the system design for better performance and efficiency.

Instruction Level Parallelism (ILP)

Multiple instructions can be executed concurrently.

> Thread Level Parallelism (TLP)

Multiple threads can be executed concurrently.

Request Level Parallelism (RLP)

Multiple requests can be executed concurrently.

Data Level Parallelism (DLP)

Instructions operate concurrently on several data.

> Other emerging parallelisms...





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Impact of Application Heterogeneity



The application heterogeneity will cause the variation of microservice execution time under various user invocations.

- It is often caused by the actual execution logic of the microservice program.
- Different microservice exhibits various degree of application heterogeneity.

Impact of Resource Provisioning



The performance implication of microservice resource relationship is complex.

Impact of Communication Overhead



Stochastic noises in the microservice environment add up the complexity of scheduling of normal system operations.

- > Here we focus on the variation of **communication overhead** between microservices.
- > Communication time variations on single machine are more stable than across machines.

Summary of Design Challenges



The crux of the problem is two-fold:

- We need to know how different microservices should be coalesced.
- > We must ensure fairly accurate alignment throughout the process.

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Potential of Microservice Level Parallelism

Parallelism	ILP	TLP	MLP	RLP
Scheduling Level	Chip Level		Systen	n Level
Granularity	Instruction	Instruction Stream	Microservice	Monolithic Application
Key Opti. Approach	Temporal	Spatial	Temporal	Spatial
Microscopic		Critical Void	ر Macroscopic	

MLP is critical and necessary:

- > MLP is orthogonal to existing parallelism model.
- MLP focuses on microservice chain scheduling.
- > MLP considers interrelationship and uncertainty of microservices.

Core Metric: Volatility of Requests

Based on characterization, microservices exhibit volatile behaviors.

> We define volatility of request (V_r) , indicating the likelihood of the request to deviate from its ideal execution conditions.

$$V_r = \alpha \times \sum_{i=1}^n I_i \times S_i \times C_i/n$$

Abbr.	Value Range	Descriptions
I	1(low) – 3(high)	Inner logic variability
S	1(low) – 3(high)	Sensitivity to resource
С	1-3 with Var(RTT)	Communication overhead

Understanding volatility helps make prudent decisions.

- Low volatility implies the start time of microservices is more predictable and less volatile.
- High volatility implies the start time of microservices is less predictable and more volatile.

Volatility-aware Microservice Level Parallelism



- v-MLP acts as the interface layer between the request handler and the server driver.
- v-MLP aims at the efficient resource management for microservices in datacenters.

Volatility-aware Microservice Level Parallelism



- Self-organizing module considers request information to coalesce microservices.
- Self-healing module handles uncertain disturbances during executions.

Design Principles of Self-organizing Module

> Periodically refreshes the status of machines in the scheduling cluster by:

- Future remaining resource status
- Future microservice execution status

Periodically reorder the request waiting queue by:

- Volatility of the request
- Arrival time of the request
- Shortest execution time of the microservices
- SLA level of the request

Assign the microservices to machines with enough resource by volatility:

- Satisfy resource demand within 1st percentile of latency for low volatility.
- Satisfy resource demand within 50th percentile of latency for mid volatility.
- Satisfy resource demand within 99th percentile of latency for high volatility.

Design Principles of Self-healing Module

Self-healing module handles uncertain disturbances in real-time execution based on the ideal microservice pipeline produced by self-organizing module.

Delay Slot Mechanism:

- Working with waiting independent microservices (nonempty delay slot).
- Advance the execution of independent microservices to fill the resource vacancy.

Resource Stretch Mechanism:

- Working with no waiting independent microservice (empty delay slot).
- > Adjust the resource usage of executing microservices to fill the resource vacancy.



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Experiment Methodology

> Experiment Platform

Characterization Platform Configurations			
Cluster	4 worker nodes (total 24 cores) + 1 manager node		
Server	Dell R730, Intel® Xeon® E5-2620		
Memory	32GB, DDR4 for each node		
Host OS	Ubuntu 18.04.5 LTS, Docker 20.10.3		
Simulation Platform Configurations			
Server	rver Dell R740 Intel® Xeon® Gold 5218		
Host OS	Ubuntu 18.04.5 LTS, Docker 20.10.3		

Experiment Benchmarks



Experiment Methodology

Evaluated workloads

Category	Requests	
Lligh V	Compose-post in SN	
nıgri v _r	getCheapest in TT	
Mid V _r	basicSearch in TT	
	Read-home-timeline in SN	
LOW V_r	Read-user-timeline in SN	

Existing scheduling schemes

Category	Scheme	Descriptions	
Simple Scheduler	FairSched	FCFS, Allocate equal resource	
	CurSched	FCFS, Allocate by current load	
Advanced Scheduler	PartProfile	Priority, Allocate by performance profile	
	FullProfile	Priority, Allocate by overall profile	
MLP scheme	v-MLP	Our Proposal	

Evaluation Results: Effectiveness & Efficiency



Evaluation Results: Performance



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Conclusions

Insights

- Microservice characteristics
- Potential of new parallelism

Microservice Level Parallelism

- Interface layer between upper and lower
- Coordinate various microservice chains
- Tackle the uncertainty in dynamic cloud

> Help for next-generation cloud-native design.

> We will expand MLP towards more directions in future.

Thank You!

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